

A	B	C	Φ
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	0
1	0	0	0
1	0	1	0
1	1	0	1
1	1	1	1

$$\Phi = (\bar{A}\bar{B}C) + (A\bar{B}\bar{C}) + (AB\bar{C}) + (ABC)$$

$$\bar{\Phi} = (\bar{A}\bar{B}\bar{C}) + (\bar{A}B\bar{C}) + (\bar{A}BC) + (A\bar{B}C)$$

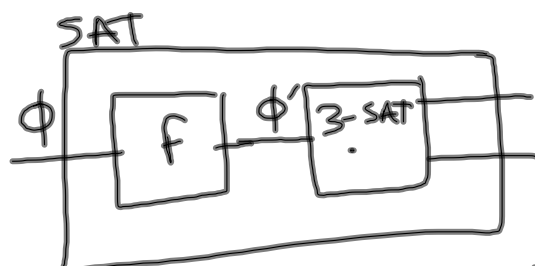
$$\Phi = \overline{(\bar{A}\bar{B}\bar{C})} \overline{(\bar{A}B\bar{C})} \overline{(\bar{A}BC)} \overline{(A\bar{B}C)}$$

$$= (A+B+C)(A+\bar{B}+C)(A+\bar{B}+\bar{C})(\bar{A}+B+\bar{C})$$

3-SAT is NPC

1. show in NP
- same as SAT
2. show $\forall L \in NP$
 $L \leq_p 3\text{-SAT}$

So show $SAT \leq_p 3\text{-SAT}$



Suppose ϕ is a Boolean formula in CNF

Construct ϕ' , a Boolean formula in 3CNF.

(term) clauses in ϕ	convert to ϕ'
3 literals	copy
1 literal x	$(x + x + x)$
2 literals $x+y$	$(x + y + y)$
4 literals $(x_1 + x_2 + x_3 + x_4)$	$(x_1 + x_2 + a) \wedge$ $(\bar{a} + x_3 + x_4)$
5 literals $(x_1 + \dots + x_5)$	$(x_1 + x_2 + a) \wedge$ $(\bar{a} + x_3 + b) \wedge$ $(\bar{b} + x_4 + x_5)$

$\therefore 3\text{ SAT is NPC}$

CLIQUE is NPC

1. done
2. $3\text{SAT} \leq_p \text{CLIQUE}$

HAMPATH = $\{ \langle G, s, t \rangle \mid \text{directed graph } G \text{ has a Hamiltonian path from } s \text{ to } t \}$

1. Show in NP \checkmark
2. $\exists \text{SAT} \leq_p \text{HAMPATH}$

$$\Phi = (a_1 \vee b_1 \vee c_1) \wedge (a_2 \vee b_2 \vee c_2) \wedge \dots \wedge (a_k \vee b_k \vee c_k)$$

variables: x_1, \dots, x_l

1. add nodes to G for each clause C_1, \dots, C_k

for each var.

$\exists k+1$ nodes

