

$$A = \{ w w w \mid w \in \{a, b\}^* \}$$

Suppose A is regular w/ pumping length p .

~~$$a^p b a^p \notin A$$~~

Consider $s = a^p b a^p b a^p b^p$

now $s \in A$ and $|s| > p$ so s can be split into 3 pieces
 $s = xyz$ s.t. $xy^iz \in A$
 for $i \geq 0$.

Consider the cases:

y is all a 's
 $xyy^iz \notin A$

one section has more a 's
 since $|xy| < p$, this is the only case.

$$s = a^p a^p a^p \quad s \in A \quad \checkmark$$

eg. $a a a a a a a a a$
 $\underbrace{\hspace{1.5cm}}_y$
 $\underbrace{a a a}_{y} \underbrace{a a a}_{y} a a a a a \in A$

$$B = \{ a^{2^n} \mid n \geq 0 \} \quad \Sigma = \{ a \}$$

$$S = a^{2^p}$$

y is at most p a 's

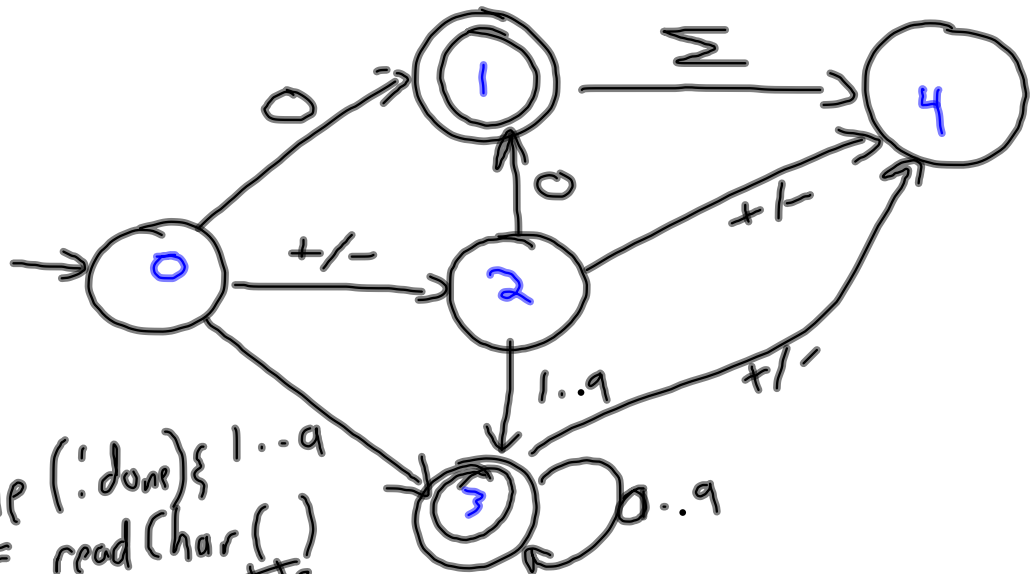
$$xyyz \stackrel{\text{is at most}}{=} a^{2^p + p} < a^{2^p + 2^p} = a^{2^{(p+1)}}$$

so $xyyz \notin B$

DFA

Finite State Machines

$$\Sigma = \{+, -, 0..9\}$$



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while (!done) {
  c = read(char);
  switch (state) {
    case 0: if (c == '0')
             state = 1;
  }
}

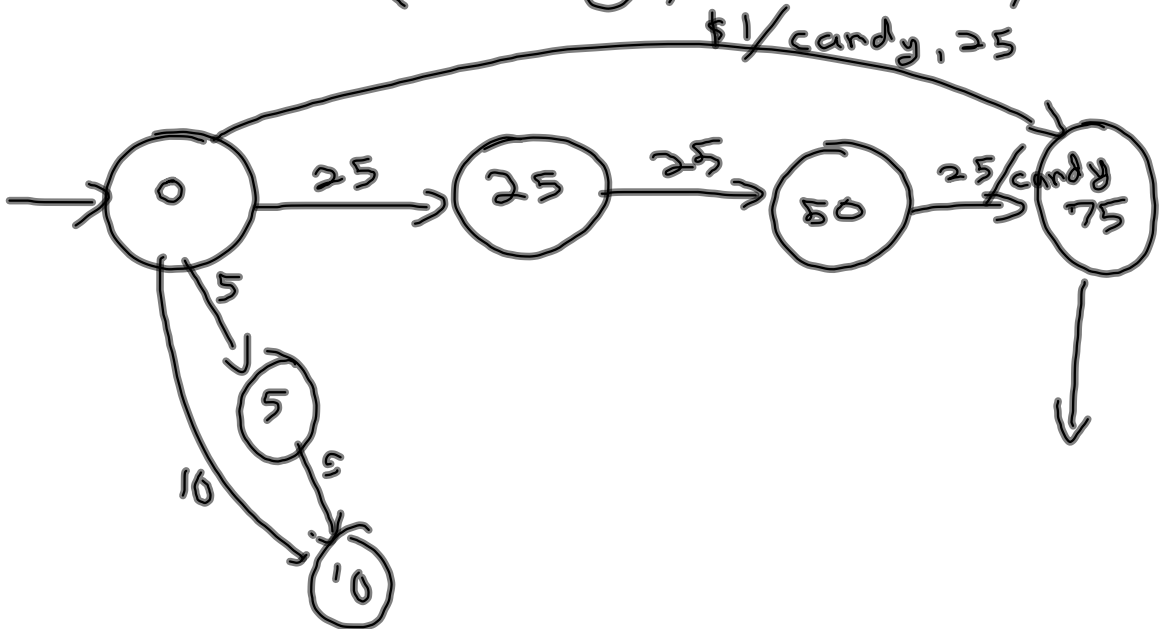
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break;

Vending

$$\Sigma = \{25, 5, 10, \$1\}$$

$$\Gamma = \{\text{candy}, 25, 5, 10\}$$



$$\{w \mid w = w^R\}$$

$$\Sigma = \{0, 1\}$$

$$\begin{array}{l} 01110 \\ 101101 \\ 010 \end{array} \}$$

$$w w^R$$

$$w \Sigma w^R$$

$$S = 0^p 1 0^p$$