

$r_i(x)$

conflicts

- access same data
- different transactions
- one must be a write

Correct Execution

- serial schedule (no interleaving)

- serializable schedule is equiv.
to some serial sched.

Equivalent?

- same results

- conflict equiv.

- order of two conflicting
ops. is the same in
both sched.

Testing for conflict serializability (Algorithm 21.1):

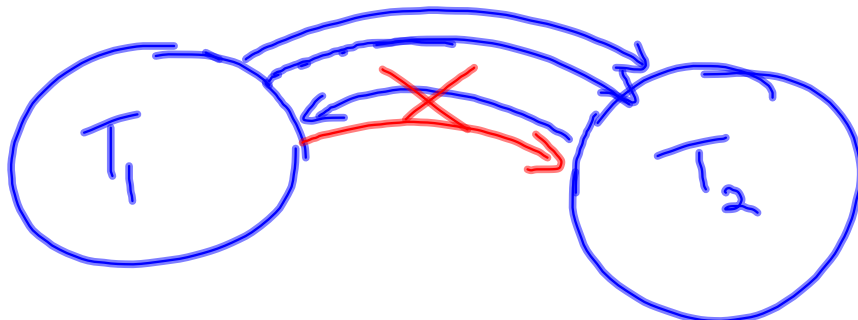
1. Create a directed graph with a node for each transaction T participating in S .
2. For each conflict between an operation in T_i and an operation in T_j , add a directed edge from T_i to T_j if the conflicting operation comes first in T_i , otherwise add an edge from T_j to T_i .
3. The schedule is serializable iff the graph has no cycles.

$$\overline{T}_1: r_1(x) w_1(x)$$

$$\overline{T}_2: r_2(y) w_2(y) r_2(x) w_2(x)$$

$$S_1: r_1(x) r_2(y) w_2(y) r_2(x) w_1(x) w_2(x)$$

conflicts	
$r_1(x)$	$w_2(x)$
$r_2(x)$	$w_1(x)$
$w_1(x)$	$w_2(x)$



$$S_1': r_1(x) r_2(y) w_2(y) \underline{w_1(x)} r_2(x) w_2(x)$$

Testing for conflict serializability (Algorithm 21.1):

1. Create a directed graph with a node for each transaction T participating in S .
2. For each conflict between an operation in T_i and an operation in T_j , add a directed edge from T_i to T_j if the conflicting operation comes first in S , otherwise add an edge from T_j to T_i .
3. The schedule is serializable iff the graph has no cycles.

Transactions:

T_1 : $r_1(X)$; $r_1(Z)$; $w_1(X)$;

T_2 : $r_2(Z)$; $r_2(Y)$; $w_2(Z)$; $w_2(Y)$;

T_3 : $r_3(X)$; $r_3(Y)$; $w_3(Y)$;

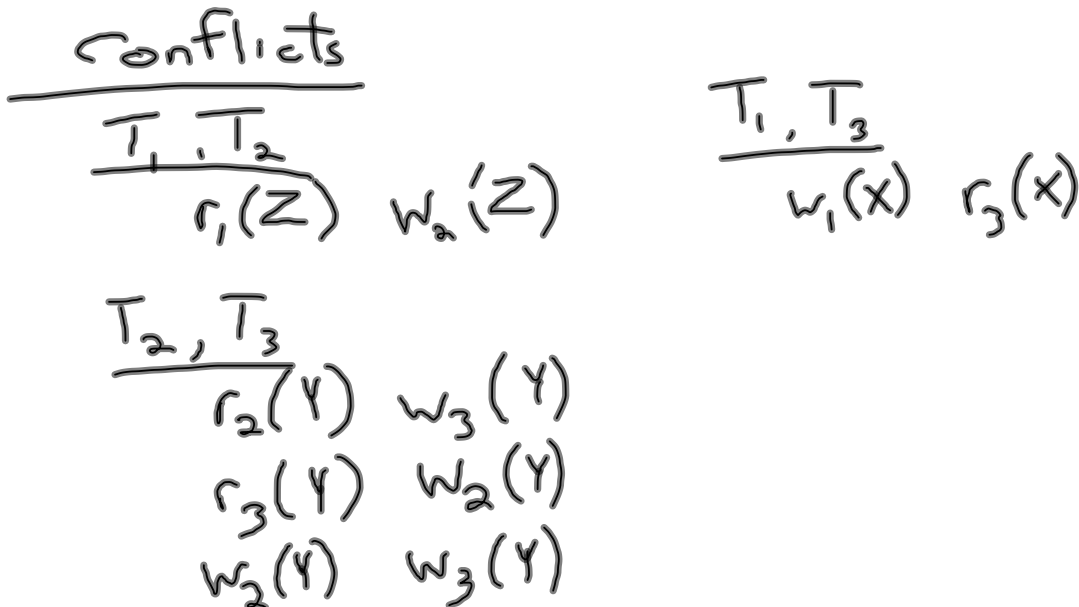
Schedules:

S_1 : $r_1(X)$; $r_2(Z)$; $r_1(Z)$; $r_3(X)$; $r_3(Y)$; $w_1(X)$; $w_3(Y)$; $r_2(Y)$; $w_2(Z)$; $w_2(Y)$;

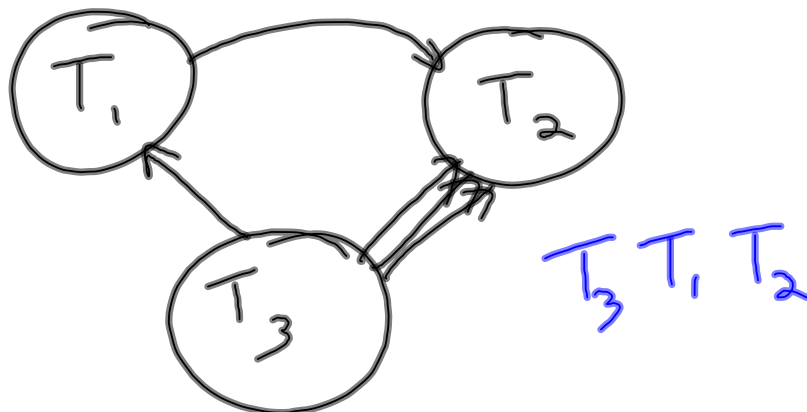
S_2 : $r_1(X)$; $r_2(Z)$; $r_3(X)$; $r_1(Z)$; $r_2(Y)$; $r_3(Y)$; $w_1(X)$; $w_2(Z)$; $w_3(Y)$; $w_2(Y)$;

S_3 : $r_2(Z)$; $r_1(X)$; $r_2(Y)$; $w_2(Z)$; $r_1(Z)$; $w_1(X)$; $r_3(X)$; $w_2(Y)$; $r_3(Y)$; $w_3(Y)$;

1. Identify the conflicts between all transactions.
2. Determine if each schedule is serializable.
3. If the schedule is serializable, what is an equivalent serial schedule?



S_1 : $r_1(X)$; $r_2(Z)$; $r_1(Z)$; $r_3(X)$; $r_3(Y)$; $w_1(X)$; $w_3(Y)$; $r_2(Y)$; $w_2(Z)$; $w_2(Y)$;



SQL

{ START TRANSACTION
{ BEGIN

COMMIT

ROLLBACK