

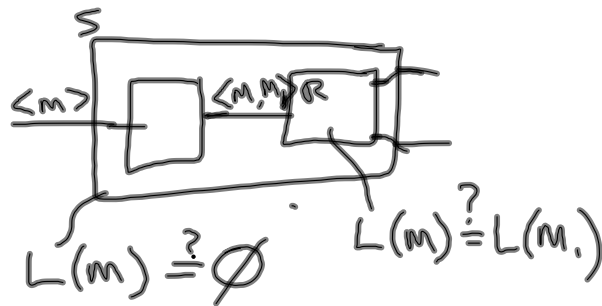
$$EQ_{TM} = \left\{ \langle M_1, M_2 \rangle \mid \begin{array}{l} M_1, M_2 \text{ are} \\ \text{TM and } L(M_1) = L(M_2) \end{array} \right\}$$

EQ_{TM} is undecidable

proof by contradiction:

assume EQ_{TM} is decidable
by TM, R .

Show that E_{TM} is decidable
(which is a contradiction)



Construct S : on input $\langle m \rangle$

1. Run R on input $\langle M, M \rangle$

where M is a TM that
rejects all input.

2. If R accepts, accept.
If R rejects, reject.

So S decides E_{TM} which
is a contradiction.

Therefore there is no decider
for EQ_{TM} .

configurations

10110q₃10110

computational history
sequence of configs

$C_1, C_2, C_3, \dots, C_\ell$

Accepting comp. history

- a) C_1 is start config.
- b) C_ℓ is an accept config
- c) each C_i follows from C_{i-1} according to rules for M .

Linear Bounded Automata (LBA)

- like a TM
- tape is only as long as the input.

LBA's accept CSLs

can decide $A_{DFA}, A_{CFG},$
 E_{DFA}, E_{CFG}

M is a LBA

q states

g symbols

n input length

There are exactly $q \cdot n \cdot g^n$
distinct configurations

A_{LBA} is decidable

E_{LBA} is undecidable

$$10q_2101 \rightarrow 100q_301 \rightarrow 10q_2101$$

Post Correspondence Problem (PCP)

given

$$\left\{ \begin{bmatrix} b \\ ca \end{bmatrix}, \begin{bmatrix} a \\ ab \end{bmatrix}, \begin{bmatrix} ca \\ a \end{bmatrix}, \begin{bmatrix} abc \\ c \end{bmatrix} \right\}$$

$$\begin{bmatrix} a \\ ab \end{bmatrix} \begin{bmatrix} b \\ ca \end{bmatrix} \begin{bmatrix} ca \\ a \end{bmatrix} \begin{bmatrix} a \\ ab \end{bmatrix} \begin{bmatrix} abc \\ c \end{bmatrix}$$

abca^aabc