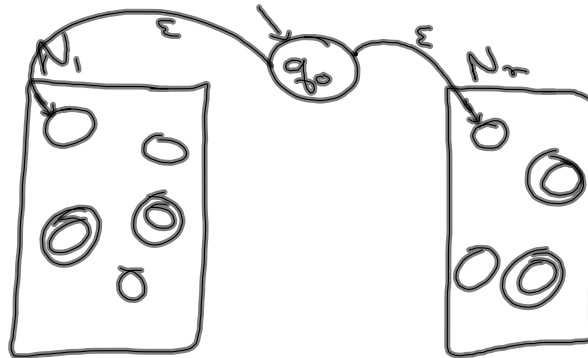


if  $A_1$  and  $A_2$  are regular  
then  $A_1 \cup A_2$  is regular.

$N_1$  is an NFA recognizing  $A_1$   
 $N_2$  " " " " "  $A_2$



$$N_1 (Q_1, \Sigma, \delta_1, q_1, F_1)$$

$$N_2 (Q_2, \Sigma, \delta_2, q_2, F_2)$$

construct

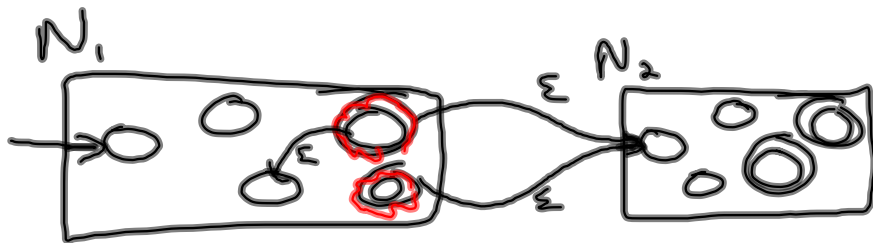
$$N (Q, \Sigma, \delta, q_0, F)$$

$$Q = Q_1 \cup Q_2 \cup \{q_0\}$$

$$\delta(q, a) = \begin{cases} \delta_1(q, a) & q \in Q_1 \\ \delta_2(q, a) & q \in Q_2 \\ \{q_1, q_2\} & q = q_0 \text{ and } a = \epsilon \\ \{\} & q = q_0 \text{ and } a \neq \epsilon \end{cases}$$

$$F = F_1 \cup F_2$$

If  $A_1$  and  $A_2$  are regular  
 then  $A_1 \circ A_2$  is regular



$A_1 = \{0, 1\}$   
 $A_2 = \{A, B\}$   
 $A_1 \circ A_2 = \{0B, 1B, 0A, 1A\}$

$N_1(Q_1, \Sigma, \delta_1, q_0, F_1)$

$N_2(Q_2, \Sigma, \delta_2, q_0, F_2)$

$N(Q, \Sigma, \delta, q_0, F_2)$

$$Q = Q_1 \cup Q_2$$

$$\delta(q, a) =$$

$$\begin{cases} \delta_1(q, a) \\ \delta_2(q, a) \\ \delta_1(q, a) \end{cases}$$

$q \in Q$   
 and  $q \notin F_1$   
 $q \in Q_2$

$q \in F_1$   
 and  $a \neq \epsilon$

$$\{q_0\} \cup \delta_1(q, a) \quad q \in F_1 \text{ and } a = \epsilon$$

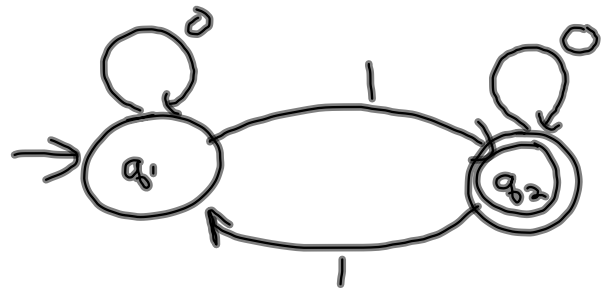
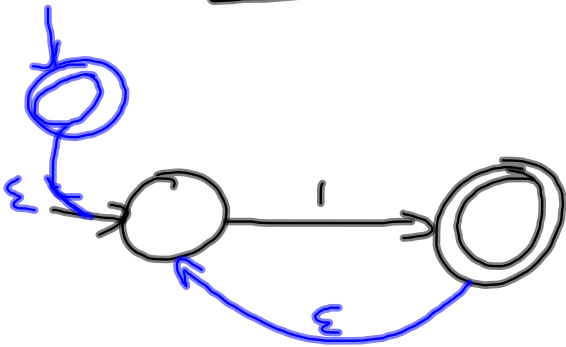
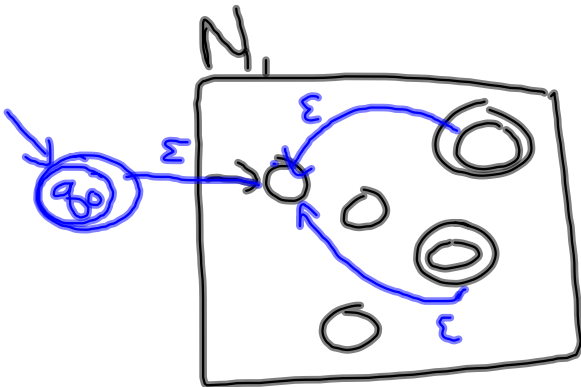
$A^*$

$$A = \{0, 1\}$$

if  $A$  is regular  
then  $A^*$  is regular

$$A^* = \{ \epsilon, 0, 1, 00, 01, 10, 11, 000, 001, 010, \dots \}$$

$$A^+ = A \circ A^*$$



$$A = \{ \text{odd # of } 1\text{'s} \}$$

$$A^* = \{ \epsilon, 1, 11, 111, \dots \}$$

$$00 \notin A^*$$