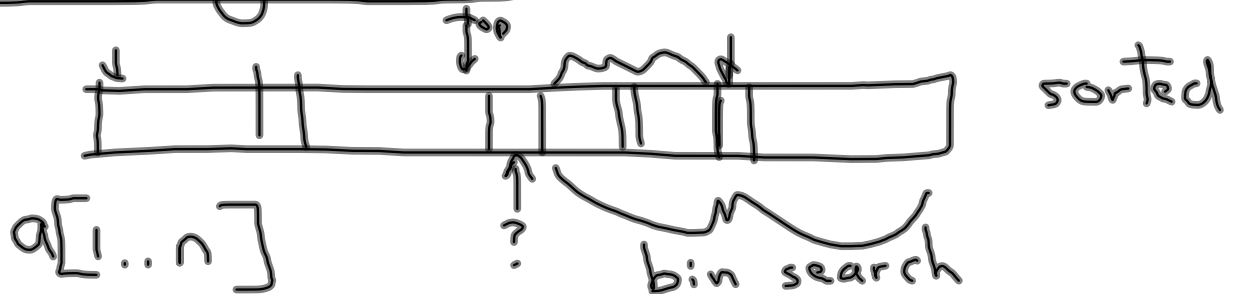


# Binary Search



index = 0

bottom = 1

top = n

while (top  $\geq$  bottom && index == 0) {

mid =  $\lfloor \text{bottom} + \text{top} / 2 \rfloor$ ;

if (a[mid] == x)

index = mid;

else if (a[mid] > x)

top = mid - 1;

else

bottom = mid + 1;

}

$w_n$ : # of iterations of the loop  
in worst case execution  
of binary search

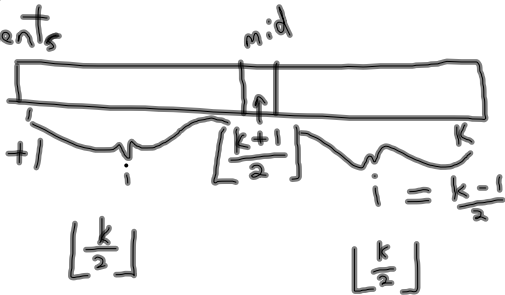
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$$w_1 = 1$$

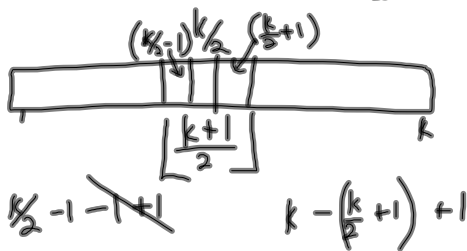
$k$  - elements

$k$  is odd

$$k = 2 \cdot i + 1$$



$k$  is even



$$w_i = 1$$

$$w_k = 1 + w_{\lfloor \frac{k}{2} \rfloor}$$

- $w_2 = 1 + w_1 = 2$
- $w_3 = 1 + w_1 = 2$
- $w_4 = 1 + w_2 = 3$
- $w_5 = 1 + w_2 = 3$
- $w_6 = 1 + w_3 = 3$
- $w_7 = 1 + w_3 = 3$
- $w_8 = 1 + w_4 = 4$
- $w_9 = 1 + w_4 = 4$
- $w_{10} = 1 + w_5 = 4$
- ...
- $w_{15} = 1 + w_7 = 4$
- $w_{16} = 1 + w_8 = 5$
- ...

if  $2^i \leq n < 2^{i+1}$   
 $w_i = i + 1$   ~~$\neq$~~

---

$k \in \mathbb{Z}, x \in \mathbb{R}$   
 and  $2^k \leq x < 2^{k+1}$   
 then  $\lfloor \log_2 x \rfloor = k$   
 (9.4.2)

$$w_n = \lfloor \log_2 n \rfloor + 1$$

$$w_1 = 1$$

$$w_n = 1 + w_{\lfloor n/2 \rfloor} \text{ for } n > 1$$

prove  $w_n = \lfloor \log_2 n \rfloor + 1 \quad n \geq 1$

Strong induction:

Base case: suppose  $n = 1$

$$w_1 = \lfloor \log_2 1 \rfloor + 1 = \lfloor 0 \rfloor + 1 = 1$$

Induction:

for all  $k \geq 2$  if

$$w_i = \lfloor \log_2 i \rfloor + 1 \text{ for}$$

$1 \leq i < k$  then

$$w_k = \lfloor \log_2 k \rfloor + 1$$

case 1:  $k$  is odd

$$\lfloor \frac{k}{2} \rfloor = \frac{k-1}{2}$$

$$w_k = 1 + w_{\lfloor k/2 \rfloor} \quad \text{def of } w$$

$$= 1 + w_{\frac{k-1}{2}}$$
$$= 1 + \left( \lfloor \log_2 \left( \frac{k-1}{2} \right) \rfloor + 1 \right)$$

$$= \lfloor \log_2(k-1) - \log_2(2) \rfloor + 2$$

$$= \lfloor \log_2(k-1) \rfloor - 1 + 2$$

for any odd int.  $k > 1$   
 $\lfloor \log_2(k-1) \rfloor = \lfloor \log_2 k \rfloor$

$$= \lfloor \log_2(k) \rfloor + 1$$

~~w.o. h.o.~~  
QED